

## 7 Two Versions of Capacity

In digital transmission systems, bits are represented by the analog signals  $s_0(t)$  and  $s_1(t)$ , both having duration  $T$ . When this signaling strategy is used over an additive white Gaussian noise channel, the probability of error in decoding the transmission of equally likely bits is given by

$$P_e = Q \left( \sqrt{\frac{\|s_1 - s_0\|^2}{2N_0}} \right),$$

where  $\|s\|^2 = \int_0^T s^2(t) dt \equiv E$ , the signal's energy, and  $Q(x)$  is the tail-integral of the standardized Gaussian probability density function.

$$Q(x) = \int_x^\infty \frac{1}{\sqrt{2\pi}} e^{-\alpha^2/2} d\alpha$$

Imposing an energy constraint on the two signals,  $E_{s_0}, E_{s_1} \leq E_0$ , the smallest achievable probability of error results when  $s_1(t) = -s_0(t)$  and each signal has the largest possible energy. From a discrete-channel viewpoint, this encoding-analog transmission-decoding implementation can be accurately modeled as a binary symmetric channel (BSC) with a cross-over probability equal to  $P_e$ .

- (a) Find an expression for and plot the capacity of this channel as a function of the signal-to-noise ratio.

We also have an expression for the capacity of an additive white Gaussian noise channel having bandwidth  $W$ .

$$C = W \log \left( 1 + \frac{P}{N_0 W} \right) \text{ bits/s}$$

Here,  $P$  is the signal's power. The issue becomes reconciling the two ways of computing capacity.

- (b) Before performing any calculations, which way of computing capacity yields the larger answer or are they the same?
- (c) Calculate expressions for the two capacities so that they can be compared. Plot them to determine how they vary with signal-to-noise ratio.