

# Model reduction of large-scale dynamical systems

## Lecture III: Krylov approximation and rational interpolation

**Thanos Antoulas**

**Rice University and Jacobs University**

email: [aca@rice.edu](mailto:aca@rice.edu)

URL: [www.ece.rice.edu/~aca](http://www.ece.rice.edu/~aca)

International School, Monopoli, 7 - 12 September 2008

# Outline

- 1 Krylov approximation methods
- 2 The Arnoldi and the Lanczos procedures
  - The Arnoldi procedure
  - The Lanczos procedure
  - An example
- 3 Krylov methods and moment matching
  - Remarks
- 4 Rational interpolation by Krylov projection
  - Realization by projection
  - Interpolation by projection
- 5 Choice of Krylov projection points: Optimal  $\mathcal{H}_2$  model reduction
- 6 Summary: Lectures II and III

# Outline

- 1 Krylov approximation methods
- 2 The Arnoldi and the Lanczos procedures
  - The Arnoldi procedure
  - The Lanczos procedure
  - An example
- 3 Krylov methods and moment matching
  - Remarks
- 4 Rational interpolation by Krylov projection
  - Realization by projection
  - Interpolation by projection
- 5 Choice of Krylov projection points: Optimal  $\mathcal{H}_2$  model reduction
- 6 Summary: Lectures II and III

## Krylov approximation methods

Given  $\Sigma = \left( \begin{array}{c|c} \mathbf{A} & \mathbf{B} \\ \hline \mathbf{C} & \mathbf{D} \end{array} \right)$ , expand the transfer function around  $s_0$ :

$$\mathbf{H}(s) = \eta_0 + \eta_1(s - s_0) + \eta_2(s - s_0)^2 + \eta_3(s - s_0)^3 + \dots$$

**Moments at  $s_0$ :**  $\eta_j, j \geq 0$ . Find  $\hat{\Sigma} = \left( \begin{array}{c|c} \hat{\mathbf{A}} & \hat{\mathbf{B}} \\ \hline \hat{\mathbf{C}} & \hat{\mathbf{D}} \end{array} \right)$ , with

$$\hat{\mathbf{H}}(s) = \hat{\eta}_0 + \hat{\eta}_1(s - s_0) + \hat{\eta}_2(s - s_0)^2 + \hat{\eta}_3(s - s_0)^3 + \dots$$

such that for appropriate  $k$ :

$$\eta_j = \hat{\eta}_j, \quad j = 1, 2, \dots, k$$

**Moment matching** methods can be implemented in a **numerically stable** and **efficient way**.

## Krylov approximation methods: Special cases

- $s_0 = \infty$  *Moments: Markov parameters*

*Problem: (partial) realization*

*Solution computed through: Lanczos and Arnoldi procedures*

- $s_0 = 0$

*Problem: Padé approximation*

*Solution computed through: Lanczos and Arnoldi procedures*

- *In general: arbitrary  $s_0 \in \mathbb{C}$*

*Problem: Rational interpolation*

*Solution computed through: Rational Lanczos*

- **Computation of moments:** numerically problematic
- **Key fact for numerical reliability:** If  $(\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$  given
- **moment matching without moment computation**  
 $\Rightarrow$  **iterative implementation.**

# Outline

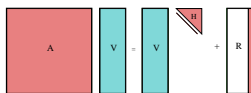
- 1 Krylov approximation methods
- 2 The Arnoldi and the Lanczos procedures
  - The Arnoldi procedure
  - The Lanczos procedure
  - An example
- 3 Krylov methods and moment matching
  - Remarks
- 4 Rational interpolation by Krylov projection
  - Realization by projection
  - Interpolation by projection
- 5 Choice of Krylov projection points: Optimal  $\mathcal{H}_2$  model reduction
- 6 Summary: Lectures II and III

## The Arnoldi procedure

Given is  $\mathbf{A} \in \mathbb{R}^{n \times n}$ , and  $\mathbf{b} \in \mathbb{R}^n$ . Let  $\mathcal{R}_k(\mathbf{A}, \mathbf{b}) \in \mathbb{R}^{n \times k}$  be the reachability or Krylov matrix. It is *assumed* that  $\mathcal{R}_k$  has full column rank equal to  $k$ .

Devise a process which is iterative and at the  $k^{\text{th}}$  step we have

$$\mathbf{A}\mathbf{V}_k = \mathbf{V}_k\mathbf{H}_k + \mathbf{R}_k, \quad \mathbf{V}_k, \mathbf{R}_k \in \mathbb{R}^{n \times k}, \quad \mathbf{H}_k \in \mathbb{R}^{k \times k}, \quad k = 1, 2, \dots, n$$



These quantities have to satisfy the following conditions at each step.

- The columns of  $\mathbf{V}_k$  are orthonormal:  $\mathbf{V}_k^* \mathbf{V}_k = \mathbf{I}_k$ ,  $k = 1, 2, \dots, n$ .
- $\text{span col } \mathbf{V}_k = \text{span col } \mathcal{R}_k(\mathbf{A}, \mathbf{b})$ ,  $k = 1, 2, \dots, n$
- The residual  $\mathbf{R}_k$  satisfies the Galerkin condition:  $\mathbf{V}_k^* \mathbf{R}_k = \mathbf{0}$ ,  $k = 1, 2, \dots, n$ .

This problem can be solved by the **Arnoldi procedure**.

## Arnoldi: recursive implementation

**Given:**  $\mathbf{A} \in \mathbb{R}^{n \times n}$ ,  $\mathbf{b} \in \mathbb{R}^n$

**Find:**  $\mathbf{V} \in \mathbb{R}^{n \times k}$ ,  $\mathbf{f} \in \mathbb{R}^n$ , and  $\mathbf{H} \in \mathbb{R}^{k \times k}$ , such that

$$\mathbf{AV} = \mathbf{VH} + \mathbf{f}\mathbf{e}_k^* \quad \text{where} \quad \mathbf{H} = \mathbf{V}^*\mathbf{AV}, \quad \mathbf{V}^*\mathbf{V} = \mathbf{I}_k, \quad \mathbf{V}^*\mathbf{f} = 0,$$

with  $\mathbf{H}$  in *upper Hessenberg* form.

- ①  $\mathbf{v}_1 = \frac{\mathbf{b}}{\|\mathbf{b}\|}$ ,  $\mathbf{w} = \mathbf{Av}_1$ ;  $\alpha_1 = \mathbf{v}_1^*\mathbf{w}$   
 $\mathbf{f}_1 = \mathbf{w} - \mathbf{v}_1\alpha_1$ ;  $\mathbf{V}_1 = (\mathbf{v}_1)$ ;  $\mathbf{H}_1 = (\alpha_1)$
- ② For  $j = 1, 2, \dots, k-1$ 
  - ①  $\beta_j = \|\mathbf{f}_j\|$ ,  $\mathbf{v}_{j+1} = \frac{\mathbf{f}_j}{\beta_j}$
  - ②  $\mathbf{V}_{j+1} = (\mathbf{V}_j \ \mathbf{v}_{j+1})$ ,  $\hat{\mathbf{H}}_j = \begin{pmatrix} \mathbf{H}_j \\ \beta_j \mathbf{e}_j^* \end{pmatrix}$
  - ③  $\mathbf{w} = \mathbf{Av}_{j+1}$ ,  $\mathbf{h} = \mathbf{V}_{j+1}^*\mathbf{w}$ ,  $\mathbf{f}_{j+1} = \mathbf{w} - \mathbf{V}_{j+1}\mathbf{h}$
  - ④  $\mathbf{H}_{j+1} = \begin{pmatrix} \hat{\mathbf{H}}_j & \mathbf{h} \end{pmatrix}$



## Properties of Arnoldi

- $\mathbf{H}_k$  is obtained by projecting  $\mathbf{A}$  onto the span of the columns of  $\mathbf{V}_k$ :  $\mathbf{H}_k = \mathbf{V}_k^* \mathbf{A} \mathbf{V}_k$ .
- The remainder  $\mathbf{R}_k$  has rank one and can be written as  $\mathbf{R}_k = \mathbf{r}_k \mathbf{e}_k^*$ , where  $\mathbf{e}_k$  is the  $k^{\text{th}}$  unit vector; thus  $\mathbf{r}_k \perp \mathcal{R}_k$ .
- This further implies that  $\mathbf{v}_{k+1} = \frac{\mathbf{r}_k}{\|\mathbf{r}_k\|}$ , where  $\mathbf{v}_{k+1}$  is the  $(k+1)^{\text{st}}$  column of  $\mathbf{V}$ . Consequently,  $\mathbf{H}_k$  is an **upper Hessenberg** matrix.

$$\mathbf{H}_k = \begin{bmatrix} h_{1,1} & h_{1,2} & h_{1,3} & \cdots & h_{1,k-1} & h_{1,k} \\ h_{2,1} & h_{2,2} & h_{2,3} & \cdots & h_{2,k-1} & h_{2,k} \\ & h_{3,2} & h_{3,3} & & h_{3,k-1} & h_{3,k} \\ & & & \ddots & \vdots & \vdots \\ & & & & h_{k-1,k-1} & h_{k-1,k} \\ & & & & h_{k,k-1} & h_{k,k} \end{bmatrix}$$

- Let  $\mathbf{p}_k(\lambda) = \det(\lambda \mathbf{I}_k - \mathbf{H}_k)$ , be the characteristic polynomial of  $\mathbf{H}_k$ . This monic polynomial is the solution of the following minimization problem

$$\mathbf{p}_k = \arg \min \|\mathbf{p}(\mathbf{A})\mathbf{b}\|_2$$

where the minimum is taken over all *monic* polynomials  $\mathbf{p}$  of degree  $k$ . Since  $\mathbf{p}_k(\mathbf{A})\mathbf{b} = \mathbf{A}^k \mathbf{b} + \mathcal{R}_k \cdot \underline{\mathbf{p}}$ , where  $\underline{\mathbf{p}}_{i+1}$  is the coefficient of  $\lambda^i$  of the polynomial  $\mathbf{p}_k$ , it also follows that the coefficients of  $\mathbf{p}_k$  provide the least squares fit between  $\mathbf{A}^k \mathbf{b}$  and the columns of  $\mathcal{R}_k$ .

- There holds

$$\mathbf{r}_k = \frac{1}{\|\mathbf{p}_{k-1}(\mathbf{A})\mathbf{b}\|} \mathbf{p}_k(\mathbf{A})\mathbf{b}, \quad \mathbf{H}_{k,k-1} = \frac{\|\mathbf{p}_k(\mathbf{A})\mathbf{b}\|}{\|\mathbf{p}_{k-1}(\mathbf{A})\mathbf{b}\|}$$

## An alternative way of looking at Arnoldi

Consider a matrix  $\mathbf{A} \in \mathbb{R}^{n \times n}$ , a starting vector  $\mathbf{b} \in \mathbb{R}^n$ , and the corresponding reachability matrix  $\mathcal{R}_n = [\mathbf{b} \ \mathbf{A}\mathbf{b} \ \dots \ \mathbf{A}^{n-1}\mathbf{b}]$ . The following relationship holds true:

$$\mathbf{A}\mathcal{R}_n = \mathcal{R}_n\mathbf{F} \quad \text{where} \quad \mathbf{F} = \begin{pmatrix} 0 & 0 & \dots & 0 & -\alpha_0 \\ 1 & 0 & \dots & 0 & -\alpha_1 \\ 0 & 1 & \dots & 0 & -\alpha_2 \\ & & \vdots & & \\ 0 & 0 & \dots & 1 & -\alpha_{n-1} \end{pmatrix}$$

and  $\chi_{\mathbf{A}}(s) = s^n + \alpha_{n-1}s^{n-1} + \dots + \alpha_1s + \alpha_0$ , is the characteristic polynomial of  $\mathbf{A}$ . Compute the QR factorization of  $\mathcal{R}_n$ :

$$\mathcal{R}_n = \mathbf{V}\mathbf{U}, \quad \mathbf{V}^*\mathbf{V} = \mathbf{I}_n, \quad \mathbf{U} \text{ upper triangular}$$

It follows that

$$\mathbf{AVU} = \mathbf{VUF} \Rightarrow \mathbf{AV} = \mathbf{V} \underbrace{\mathbf{UFU}^{-1}}_{\bar{\mathbf{A}}} \Rightarrow \mathbf{AV} = \mathbf{V}\bar{\mathbf{A}}$$

Since  $\mathbf{U}$  is upper triangular, so is  $\mathbf{U}^{-1}$ ; furthermore  $\mathbf{F}$  is upper Hessenberg. Therefore  $\bar{\mathbf{A}}$  being the product of an upper triangular times an upper Hessenberg times an upper triangular matrix is *upper Hessenberg*.

The  $k$ -step Arnoldi factorization can now be obtained by considering the first  $k$  columns of the above relationship, to wit:

$$[\mathbf{AV}]_k = [\mathbf{V}\bar{\mathbf{A}}]_k \Rightarrow \mathbf{A}[\mathbf{V}]_k = [\mathbf{V}]_k \bar{\mathbf{A}}_{kk} + \mathbf{f}_k^*$$

where  $\mathbf{f}$  is a multiple of the  $(k+1)$ -st column of  $\mathbf{V}$ . Notice that  $\bar{\mathbf{A}}_{kk}$  is still upper Hessenberg, while the columns of  $[\mathbf{V}]_k$  provide an orthonormal basis for the space spanned by the first  $k$  columns of the reachability matrix  $\mathcal{R}_n$ .



## Two-sided Lanczos

**The two-sided Lanczos procedure.** Given  $\mathbf{A} \in \mathbb{R}^{n \times n}$  which is not symmetric, and two vectors  $\mathbf{b}, \mathbf{c}^* \in \mathbb{R}^n$ , devise a process which is iterative and the  $k^{\text{th}}$  step there holds:

$$\mathbf{A}\mathbf{V}_k = \mathbf{V}_k\mathbf{H}_k + \mathbf{R}_k, \quad \mathbf{A}^*\mathbf{W}_k = \mathbf{W}_k\mathbf{H}_k + \mathbf{S}_k, \quad k = 1, 2, \dots, n.$$

- Biorthogonality:  $\mathbf{W}_k^*\mathbf{V}_k = \mathbf{I}_k$ ,
- span col  $\mathbf{V}_k = \text{span col } \mathcal{R}_k(\mathbf{A}, \mathbf{b})$ , span col  $\mathbf{W}_k = \text{span col } \mathcal{R}_k(\mathbf{A}^*, \mathbf{c}^*)$ ,
- Galerkin conditions:  $\mathbf{V}_k^*\mathbf{S}_k = 0$ ,  $\mathbf{W}_k^*\mathbf{R}_k = 0$ ,  $k = 1, 2, \dots, n$ .

**Remarks.**

- The second condition of the second item above can also be expressed as span rows  $\mathbf{W}_k^* = \text{span rows } \mathcal{O}_k(\mathbf{c}, \mathbf{A})$ , where  $\mathcal{O}_k$  is the observability matrix of the pair  $(\mathbf{c}, \mathbf{A})$ .
- The *assumption* for the solvability of this problem is  $\det \mathcal{O}_k(\mathbf{c}, \mathbf{A})\mathcal{R}_k(\mathbf{A}, \mathbf{b}) \neq 0$ ,  $k = 1, 2, \dots, n$ .
- The associated *Lanczos polynomials* are defined as  $\mathbf{p}_k(\lambda) = \det(\lambda\mathbf{I}_k - \mathbf{H}_k)$ , and the induced inner product is defined as  $\langle \mathbf{p}(\lambda), \mathbf{q}(\lambda) \rangle = \langle \mathbf{p}(\mathbf{A}^*)\mathbf{c}^*, \mathbf{q}(\mathbf{A})\mathbf{b} \rangle = \mathbf{c}\mathbf{p}(\mathbf{A}) \cdot \mathbf{q}(\mathbf{A})\mathbf{b}$ .

## Two-sided Lanczos: recursive implementation

**Given:** the triple  $\mathbf{A} \in \mathbb{R}^{n \times n}$ ,  $\mathbf{b}, \mathbf{c}^* \in \mathbb{R}^n$

**Find:**  $\mathbf{V}, \mathbf{W} \in \mathbb{R}^{n \times k}$ ,  $\mathbf{g} \in \mathbb{R}^n$ , and  $\mathbf{H} \in \mathbb{R}^{k \times k}$ , such that

$$\mathbf{AV} = \mathbf{VH} + \mathbf{f}\mathbf{e}_k^*, \quad \mathbf{A}^*\mathbf{W} = \mathbf{WH}^* + \mathbf{g}\mathbf{e}_k^* \quad \text{where}$$

$\mathbf{H} = \mathbf{V}^*\mathbf{AW}$ ,  $\mathbf{V}^*\mathbf{W} = \mathbf{I}_k$ ,  $\mathbf{W}^*\mathbf{f} = 0$ ,  $\mathbf{V}^*\mathbf{g} = 0$ . The projections  $\pi_L$  and  $\pi_U$  above, are given by  $\mathbf{V}^*$ ,  $\mathbf{W}$ , respectively.

- 1  $\beta_1 := \sqrt{|\mathbf{b}^*\mathbf{c}^*|}$ ,  $\gamma_1 := \text{sgn}(\mathbf{b}^*\mathbf{c}^*)\beta_1$   
 $\mathbf{v}_1 = \mathbf{b}/\beta_1$ ,  $\mathbf{w}_1 := \mathbf{c}^*/\gamma_1$
- 2 For  $j = 1, \dots, k$ , set
  - 1  $\alpha_j = \mathbf{w}_j^*\mathbf{A}\mathbf{v}_j$
  - 2  $\mathbf{r}_j = \mathbf{A}\mathbf{v}_j - \alpha_j\mathbf{v}_j - \gamma_j\mathbf{v}_{j-1}$ ,  $\mathbf{q}_j = \mathbf{A}^*\mathbf{w}_j - \alpha_j\mathbf{w}_j - \beta_j\mathbf{w}_{j-1}$
  - 3  $\beta_{j+1} = \sqrt{|\mathbf{r}_j^*\mathbf{q}_j|}$ ,  $\gamma_{j+1} = \text{sgn}(\mathbf{r}_j^*\mathbf{q}_j)\beta_{j+1}$
  - 4  $\mathbf{v}_{j+1} = \mathbf{r}_j/\beta_{j+1}$ ,  $\mathbf{w}_{j+1} = \mathbf{q}_j/\gamma_{j+1}$

## Properties of two-sided Lanczos

- $\mathbf{H}_k$  is obtained by projecting  $A$  as follows:  $\mathbf{H}_k = \mathbf{W}_k^* \mathbf{A} \mathbf{V}_k$ .
- The remainders  $\mathbf{R}_k$ ,  $\mathbf{S}_k$  have rank one and can be written as  $\mathbf{R}_k = \mathbf{r}_k \mathbf{e}_k^*$ ,  
 $\mathbf{S}_k = \mathbf{q}_k \mathbf{e}_k^*$ .
- This further implies that  $\mathbf{v}_{k+1}$ ,  $\mathbf{w}_{k+1}$  are scaled versions of  $\mathbf{r}_k$ ,  $\mathbf{q}_k$  respectively. Consequently,  $\mathbf{H}_k$  is a **tridiagonal** matrix.
- The generalized Lanczos polynomials  $\mathbf{p}_k(\lambda) = \det(\lambda \mathbf{I}_k - \mathbf{H}_k)$ ,  
 $k = 0, 1, \dots, n-1$ ,  $\mathbf{p}_0 = 1$ , are orthogonal:  $\langle \mathbf{p}_i, \mathbf{p}_j \rangle = 0$ , for  $i \neq j$ .
- The columns of  $\mathbf{V}_k$ ,  $\mathbf{W}_k$  and the Lanczos polynomials satisfy the following three-term recurrences

$$\begin{aligned}
 \gamma_k \mathbf{v}_{k+1} &= (\mathbf{A} - \alpha_k) \mathbf{v}_k & - \beta_{k-1} \mathbf{v}_{k-1} \\
 \beta_k \mathbf{w}_{k+1} &= (\mathbf{A}^* - \alpha_k) \mathbf{w}_k & - \gamma_{k-1} \mathbf{w}_{k-1} \\
 \gamma_k \mathbf{p}_{k+1}(\lambda) &= (\lambda - \alpha_k) \mathbf{p}_k(\lambda) & - \beta_{k-1} \mathbf{p}_{k-1}(\lambda) \\
 \beta_k \mathbf{q}_{k+1}(\lambda) &= (\lambda - \alpha_k) \mathbf{q}_k(\lambda) & - \gamma_{k-1} \mathbf{q}_{k-1}(\lambda)
 \end{aligned}$$

## Example: symmetric Lanczos

Consider the following symmetric matrix:

$$\mathbf{A} = \begin{bmatrix} 2 & 1 & 2 & 1 \\ 1 & 2 & 0 & 1 \\ 2 & 0 & 2 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

With the starting vector  $\mathbf{b} = [1 \ 0 \ 0 \ 0]^*$ , we obtain



$$\mathbf{V}_1 = \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}, \mathbf{H}_1 = [2], \mathbf{R}_1 = \begin{bmatrix} 0 \\ 1 \\ 2 \\ 1 \end{bmatrix}$$

$$\mathbf{V}_2 = \begin{bmatrix} 1 & 0 \\ 0 & \frac{1}{\sqrt{6}} \\ 0 & \frac{2}{\sqrt{6}} \\ 0 & \frac{1}{\sqrt{6}} \end{bmatrix}, \mathbf{H}_2 = \begin{bmatrix} 2 & \sqrt{6} \\ \sqrt{6} & \frac{8}{3} \end{bmatrix}, \mathbf{R}_2 = \begin{bmatrix} 0 & 0 \\ 0 & \frac{1}{\sqrt{54}} \\ 0 & \frac{-1}{\sqrt{54}} \\ 0 & \frac{1}{\sqrt{54}} \end{bmatrix}$$

$$\mathbf{V}_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & \frac{1}{\sqrt{6}} & \frac{1}{\sqrt{3}} \\ 0 & \frac{2}{\sqrt{6}} & \frac{-1}{\sqrt{3}} \\ 0 & \frac{1}{\sqrt{6}} & \frac{1}{\sqrt{3}} \end{bmatrix}, \mathbf{H}_3 = \begin{bmatrix} 2 & \sqrt{6} & 0 \\ \sqrt{6} & \frac{8}{3} & \frac{1}{\sqrt{18}} \\ 0 & \frac{1}{\sqrt{18}} & \frac{4}{3} \end{bmatrix}, \mathbf{R}_3 = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & \frac{\sqrt{3}}{2} \\ 0 & 0 & 0 \\ 0 & 0 & \frac{-\sqrt{3}}{2} \end{bmatrix}$$

$$\mathbf{V}_4 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & \frac{1}{\sqrt{6}} & \frac{1}{\sqrt{3}} & \frac{1}{\sqrt{2}} \\ 0 & \frac{2}{\sqrt{6}} & \frac{-1}{\sqrt{3}} & 0 \\ 0 & \frac{1}{\sqrt{6}} & \frac{1}{\sqrt{3}} & \frac{-1}{\sqrt{2}} \end{bmatrix}, \mathbf{H}_4 = \begin{bmatrix} 2 & \sqrt{6} & 0 & 0 \\ \sqrt{6} & \frac{8}{3} & \frac{1}{\sqrt{18}} & 0 \\ 0 & \frac{1}{\sqrt{18}} & \frac{4}{3} & \frac{\sqrt{3}}{\sqrt{2}} \\ 0 & 0 & \frac{\sqrt{3}}{\sqrt{2}} & 0 \end{bmatrix}, \mathbf{R}_4 = \mathbf{0}_{4 \times 4}$$

where

$$\mathbf{A}\mathbf{V}_k = \mathbf{V}_k\mathbf{H}_k + \mathbf{R}_k, \mathbf{V}_k^*\mathbf{R}_k = 0 \Rightarrow \mathbf{H}_k = \mathbf{V}_k^*\mathbf{A}\mathbf{V}_k, k = 1, 2, 3, 4.$$

# Outline

- 1 Krylov approximation methods
- 2 The Arnoldi and the Lanczos procedures
  - The Arnoldi procedure
  - The Lanczos procedure
  - An example
- 3 Krylov methods and moment matching**
  - Remarks**
- 4 Rational interpolation by Krylov projection
  - Realization by projection
  - Interpolation by projection
- 5 Choice of Krylov projection points: Optimal  $\mathcal{H}_2$  model reduction
- 6 Summary: Lectures II and III

## Arnoldi and moment matching

The Arnoldi factorization can be used for model reduction as follows. Recall the QR factorization of the reachability matrix  $\mathcal{R}_k \in \mathbb{R}^{n \times k}$ ; a projection  $\mathbf{V}\mathbf{V}^*$  can then be attached to this factorization:

$$\mathcal{R}_k = \mathbf{V}\mathbf{U} \Rightarrow \mathbf{V} = \mathcal{R}_k\mathbf{U}^{-1}$$

where  $\mathbf{V} \in \mathbb{R}^{n \times k}$ ,  $\mathbf{V}^*\mathbf{V} = \mathbf{I}_k$ , and  $\mathbf{U}$  is upper triangular. The reduced order system is:

$$\bar{\Sigma} = \left( \begin{array}{c|c} \bar{\mathbf{A}} & \bar{\mathbf{B}} \\ \hline \bar{\mathbf{C}} & \end{array} \right) \quad \text{where} \quad \boxed{\bar{\mathbf{A}} = \mathbf{V}^*\mathbf{A}\mathbf{V}}, \quad \boxed{\bar{\mathbf{B}} = \mathbf{V}^*\mathbf{B}}, \quad \boxed{\bar{\mathbf{C}} = \mathbf{C}\mathbf{V}}$$

**Theorem.**  $\bar{\Sigma}$  as defined above satisfies the equality of the Markov parameters  $\hat{\eta}_i = \eta_i$ ,  $i = 1, \dots, k$ . Furthermore,  $\bar{\mathbf{A}}$  is in Hessenberg form, and  $\bar{\mathbf{B}}$  is a multiple of the unit vector  $\mathbf{e}_1$ .

**Proof.** First notice that since  $\mathbf{U}$  is upper triangular,  $\mathbf{v}_1 = \frac{\mathbf{B}}{\|\mathbf{B}\|}$ , and since  $\mathbf{V}^*\mathcal{R}_k = \mathbf{U}$  it follows that  $\bar{\mathbf{B}} = \mathbf{u}_1 = \|\mathbf{B}\| \mathbf{e}_1$ ; therefore  $\bar{\mathbf{B}} = \mathbf{V}^*\mathbf{B}$ .  $\mathbf{V}\mathbf{V}^*\mathbf{B} = \mathbf{V}\bar{\mathbf{B}} = \mathbf{B}$ , hence  $\bar{\mathbf{A}}\bar{\mathbf{B}} = \mathbf{V}^*\mathbf{A}\mathbf{V}\mathbf{V}^*\mathbf{B} = \mathbf{V}^*\mathbf{A}\mathbf{B}$ ; in general, since  $\mathbf{V}\mathbf{V}^*$  is a projection along the columns of  $\mathcal{R}_k$ , we have  $\mathbf{V}\mathbf{V}^*\mathcal{R}_k = \mathcal{R}_k$ ; moreover:  $\hat{\mathcal{R}}_k = \mathbf{V}^*\mathcal{R}_k$ ; hence

$$(\hat{\eta}_1 \cdots \hat{\eta}_k) = \hat{\mathbf{C}}\hat{\mathcal{R}}_k = \mathbf{C}\mathbf{V}\mathbf{V}^*\mathcal{R}_k = \mathbf{C}\mathcal{R}_k = (\eta_1 \cdots \eta_k)$$

Finally, the upper triangularity of  $\mathbf{U}$  implies that  $\mathbf{A}$  is in Hessenberg form. ■

### Remark.

Similarly, one can show that reduction by means the two-sided Lanczos procedure preserves  $2k$  Markov parameters.

## Remarks

- The number of operations is  $\mathcal{O}(k^2n)$  vs.  $\mathcal{O}(n^3)$ , which implies **efficiency**. The requirement for memory is large if  $k$  is relatively large.
- Only *matrix-vector* multiplications are required. No matrix factorizations and/or inversions. There is no need to compute the transformed  $n$ -th order model and then truncate. This eliminates **ill-conditioning**.
- **Drawbacks:**
  - Numerical issue: Arnoldi/Lanczos methods lose orthogonality. This comes from the instability of the classical Gram-Schmidt procedure. Remedy: re-orthogonalization.
  - no global error bound.
  - $\hat{\Sigma}$  tends to approximate the high frequency poles of  $\Sigma$ . Remedy: match expansions around other frequencies  $\Rightarrow$  **rational Lanczos**. ■

# Outline

- 1 Krylov approximation methods
- 2 The Arnoldi and the Lanczos procedures
  - The Arnoldi procedure
  - The Lanczos procedure
  - An example
- 3 Krylov methods and moment matching
  - Remarks
- 4 Rational interpolation by Krylov projection
  - Realization by projection
  - Interpolation by projection
- 5 Choice of Krylov projection points: Optimal  $\mathcal{H}_2$  model reduction
- 6 Summary: Lectures II and III

## Partial realization by projection

Given a system  $\Sigma = (\mathbf{A}, \mathbf{B}, \mathbf{C})$ , where  $\mathbf{A} \in \mathbb{R}^{n \times n}$  and  $\mathbf{B}, \mathbf{C}^* \in \mathbb{R}^n$ , We seek a lower dimensional model  $\hat{\Sigma} = (\hat{\mathbf{A}}, \hat{\mathbf{B}}, \hat{\mathbf{C}})$ , where  $\hat{\mathbf{A}} \in \mathbb{R}^{k \times k}$ ,  $\hat{\mathbf{B}}, \hat{\mathbf{C}}^* \in \mathbb{R}^k$ ,  $k < n$ , such that  $\hat{\Sigma}$  preserves some properties of the original system, through appropriate projection methods. In other words, we seek  $\mathbf{V} \in \mathbb{R}^{n \times k}$  and  $\mathbf{W} \in \mathbb{R}^{n \times k}$  such that  $\mathbf{W}^* \mathbf{V} = \mathbf{I}_k$ , and the reduced system is given by:

$$\hat{\mathbf{A}} = \mathbf{W}^* \mathbf{A} \mathbf{V}, \quad \hat{\mathbf{B}} = \mathbf{W}^* \mathbf{B}, \quad \hat{\mathbf{C}} = \mathbf{C} \mathbf{V}.$$

### Lemma

*With  $\mathbf{V} = [\mathbf{B}, \mathbf{A}\mathbf{B}, \dots, \mathbf{A}^{k-1}\mathbf{B}] = \mathcal{R}_k(\mathbf{A}, \mathbf{B})$  and  $\mathbf{W}$  any left inverse of  $\mathbf{V}$ ,  $\hat{\Sigma}$  is a partial realization of  $\Sigma$  and matches  $k$  Markov parameters.*

From a numerical point of view, one would not use  $\mathbf{V}$  as defined above since usually the columns of  $\mathbf{V}$  are almost linearly dependent. As it turns out any matrix whose column span is the same as that of  $\mathbf{V}$  can be used.

*Proof.*

We have  $\hat{\mathbf{C}}\hat{\mathbf{B}} = \mathbf{C}\mathbf{V}\mathbf{W}^*\mathbf{B} = \mathbf{C}\mathcal{R}_k(\mathbf{A}, \mathbf{B})\mathbf{e}_1 = \mathbf{C}\mathbf{B}$ ; furthermore

$$\hat{\mathbf{C}}\hat{\mathbf{A}}^j\hat{\mathbf{B}} = \mathbf{C}\mathcal{R}_k(\mathbf{A}, \mathbf{B})\mathbf{W}^*\mathbf{A}^j\mathcal{R}_k(\mathbf{A}, \mathbf{B})\mathbf{e}_1 = \mathbf{C}\mathcal{R}_k(\mathbf{A}, \mathbf{B})\mathbf{W}^*\mathbf{A}^j\mathbf{B} = \mathbf{C}\mathcal{R}_k(\mathbf{A}, \mathbf{B})\mathbf{e}_{j+1} = \mathbf{C}\mathbf{A}^j\mathbf{B}, \quad j = 1, \dots, k-1.$$

□

## Rational interpolation by projection

Suppose now that we are given  $k$  distinct points  $s_j \in \mathbb{C}$ .  $\mathbf{V}$  is defined as the generalized reachability matrix

$$\mathbf{V} = [(\mathbf{s}_1 \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B}, \dots, (\mathbf{s}_k \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B}],$$

and as before, let  $\mathbf{W}$  be any left inverse of  $\mathbf{V}$ . Then

### Lemma

$\hat{\Sigma}$  defined above, interpolates the transfer function of  $\Sigma$  at the  $s_j$ , that is

$$\mathbf{H}(s_j) = \mathbf{C}(s_j \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B} = \hat{\mathbf{C}}(s_j \mathbf{I}_k - \hat{\mathbf{A}})^{-1} \hat{\mathbf{B}} = \hat{\mathbf{H}}(s_j), \quad j = 1, \dots, k.$$

### Proof.

The following string of equalities leads to the desired result:

$$\begin{aligned} \hat{\mathbf{C}}(s_j \mathbf{I}_k - \hat{\mathbf{A}})^{-1} \hat{\mathbf{B}} &= \mathbf{C} \mathbf{V} (s_j \mathbf{I}_k - \mathbf{W}^* \mathbf{A} \mathbf{V})^{-1} \mathbf{W}^* \mathbf{B} \\ &= \mathbf{C} [(\mathbf{s}_1 \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B}, \dots, (\mathbf{s}_k \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B}] (\mathbf{W}^* (s_j \mathbf{I}_n - \mathbf{A}) \mathbf{V})^{-1} \mathbf{W}^* \mathbf{B} \\ &= [\mathbf{C}(\mathbf{s}_1 \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B}, \dots, \mathbf{C}(\mathbf{s}_k \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B}] ([\dots \mathbf{W}^* \mathbf{B} \dots])^{-1} \mathbf{W}^* \mathbf{B} \\ &= [\mathbf{C}(\mathbf{s}_1 \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B}, \dots, \mathbf{C}(\mathbf{s}_k \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B}] \mathbf{e}_j \\ &= \mathbf{C}(s_j \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B}. \end{aligned}$$



## Matching points with multiplicity

We now wish to match the value of the transfer function at a given point  $s_0 \in \mathbb{C}$ , together with  $k - 1$  derivatives. For this we define the *generalized reachability matrix*

$$\mathbf{V} = [ (s_0 \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B}, (s_0 \mathbf{I}_n - \mathbf{A})^{-2} \mathbf{B}, \dots, (s_0 \mathbf{I}_n - \mathbf{A})^{-k} \mathbf{B} ],$$

together with any left inverse  $\mathbf{W}$  thereof.

### Lemma

$\hat{\Sigma}$  interpolates the transfer function of  $\Sigma$  at  $s_0$ , together with  $k - 1$  derivatives at the same point,  $j = 0, 1, \dots, k - 1$ :

$$\left. \frac{(-1)^j}{j!} \frac{d^j}{ds^j} \mathbf{H}(s) \right|_{s=s_0} = \mathbf{C}(s_0 \mathbf{I}_n - \mathbf{A})^{-(j+1)} \mathbf{B} = \hat{\mathbf{C}}(s_0 \mathbf{I}_k - \hat{\mathbf{A}})^{-(j+1)} \hat{\mathbf{B}} = \left. \frac{(-1)^j}{j!} \frac{d^j}{ds^j} \hat{\mathbf{H}}(s) \right|_{s=s_0}$$

### Proof.

Let  $\mathbf{V}$  be as defined as above, and  $\mathbf{W}$  be such that  $\mathbf{W}^* \mathbf{V} = \mathbf{I}_k$ . It readily follows that the projected matrix  $s_0 \mathbf{I}_r - \hat{\mathbf{A}}$  is in companion form (expression on the left) and therefore its powers are obtained by shifting its columns to the right:

$$s_0 \mathbf{I}_k - \hat{\mathbf{A}} = \mathbf{W}^* (s_0 \mathbf{I}_n - \mathbf{A}) \mathbf{V} = [\mathbf{W}^* \mathbf{B}, \mathbf{e}_1, \dots, \mathbf{e}_{k-1}] \Rightarrow (s_0 \mathbf{I}_k - \hat{\mathbf{A}})^\ell = \underbrace{[\dots]}_{\ell-1}, \mathbf{W}^* \mathbf{B}, \mathbf{e}_1, \dots, \mathbf{e}_{k-\ell}].$$

Consequently  $[\mathbf{W}^* (s_0 \mathbf{I}_n - \mathbf{A}) \mathbf{V}]^{-\ell} \mathbf{W}^* \mathbf{B} = \mathbf{e}_\ell$ , which finally implies

$$\hat{\mathbf{C}}(s_0 \mathbf{I}_k - \hat{\mathbf{A}})^{-\ell} \hat{\mathbf{B}} = \mathbf{C} \mathbf{V} [\mathbf{W}^* (s_0 \mathbf{I}_n - \mathbf{A}) \mathbf{V}]^{-\ell} \mathbf{W}^* \mathbf{B} = \mathbf{C} \mathbf{V} \mathbf{e}_\ell = \mathbf{C}(s_0 \mathbf{I}_n - \mathbf{A})^{-\ell} \mathbf{B}, \quad \ell = 1, 2, \dots, k.$$

□

## General result: rational Krylov

A projector which is composed of any combination of the above three cases achieves matching of an appropriate number of Markov parameters and moments. Let the partial reachability matrix be

$$\mathcal{R}_k(\mathbf{A}, \mathbf{B}) = [\mathbf{B} \quad \mathbf{A}\mathbf{B} \quad \dots \quad \mathbf{A}^{k-1}\mathbf{B}],$$

and partial generalized reachability matrix be:

$$\mathcal{R}_k(\mathbf{A}, \mathbf{B}; \sigma) = [(\sigma\mathbf{I}_n - \mathbf{A})^{-1}\mathbf{B} \quad (\sigma\mathbf{I}_n - \mathbf{A})^{-2}\mathbf{B} \quad \dots \quad (\sigma\mathbf{I}_n - \mathbf{A})^{-k}\mathbf{B}].$$

### Rational Krylov

**(a)** If  $\mathbf{V}$  as defined in the above three cases is replaced by  $\bar{\mathbf{V}} = \mathbf{V}\mathbf{R}$ ,  $\mathbf{R} \in \mathbb{R}^{k \times k}$ ,  $\det \mathbf{R} \neq 0$ , and  $\mathbf{W}$  by  $\bar{\mathbf{W}} = \mathbf{R}^{-1}\mathbf{W}$ , the same matching results hold true.

**(b)** Let  $\mathbf{V}$  be such that

$$\text{span col } \mathbf{V} = \text{span col } [\mathcal{R}_k(\mathbf{A}, \mathbf{B}) \quad \mathcal{R}_{m_1}(\mathbf{A}, \mathbf{B}; \sigma_1) \quad \dots \quad \mathcal{R}_{m_\ell}(\mathbf{A}, \mathbf{B}; \sigma_\ell)],$$

and  $\mathbf{W}$  any left inverse of  $\mathbf{V}$ . The reduced system matches  $k$  Markov parameters and  $m_i$  moments at  $\sigma_i \in \mathbb{C}$ ,  $i = 1, \dots, \ell$ .

Two-sided projections: the choice of  $\mathbf{W}$ 

Let  $\mathcal{O}_k(\mathbf{C}, \mathbf{A}) \in \mathbb{R}^{k \times n}$ , be the partial observability matrix consisting of the first  $k$  rows of  $\mathcal{O}_n(\mathbf{C}, \mathbf{A}) \in \mathbb{R}^{n \times n}$ . The first case is

$$\mathbf{V} = \mathcal{R}_k(\mathbf{A}, \mathbf{B}), \quad \mathbf{W} = \underbrace{(\mathcal{O}_k(\mathbf{C}, \mathbf{A})\mathcal{R}_k(\mathbf{A}, \mathbf{B}))^{-1}}_{\mathcal{H}_k} \mathcal{O}_k(\mathbf{C}, \mathbf{A}).$$

## Lemma

Assuming that  $\det \mathcal{H}_k \neq 0$ ,  $\hat{\Sigma}$  is a partial realization of  $\Sigma$  and matches  $2k$  Markov parameters.

Given  $2k$  distinct points  $s_1, \dots, s_{2k}$ , we will make use of the following generalized reachability and observability matrices:

$$\tilde{\mathbf{V}} = [ (s_1 \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B} \quad \dots \quad (s_k \mathbf{I}_n - \mathbf{A})^{-1} \mathbf{B} ], \quad \tilde{\mathbf{W}} = [ (s_{k+1} \mathbf{I}_n - \mathbf{A}^*)^{-1} \mathbf{C}^* \quad \dots \quad (s_{2k} \mathbf{I}_n - \mathbf{A}^*)^{-1} \mathbf{C}^* ].$$

## Lemma

Assuming that  $\det \tilde{\mathbf{W}}^* \tilde{\mathbf{V}} \neq 0$ , the projected system  $\hat{\Sigma}$  where  $\mathbf{V} = \tilde{\mathbf{V}}$  and  $\mathbf{W} = \tilde{\mathbf{W}}(\tilde{\mathbf{V}}^* \tilde{\mathbf{W}})^{-1}$  interpolates the transfer function of  $\Sigma$  at the  $2k$  points  $s_j$ .

## Remarks

(a) The same procedure as above can be used to approximate *implicit systems*, i.e., systems that are given in a generalized form

$\mathbf{E}\dot{\mathbf{x}}(t) = \mathbf{A}\mathbf{x}(t) + \mathbf{B}\mathbf{u}(t)$ ,  $\mathbf{y}(t) = \mathbf{C}\mathbf{x}(t)$ , where  $\mathbf{E}$  may be singular. The reduced system is given by

$$\hat{\mathbf{E}} = \mathbf{W}^* \mathbf{E} \mathbf{V}, \quad \hat{\mathbf{A}} = \mathbf{W}^* \mathbf{A} \mathbf{V}, \quad \hat{\mathbf{B}} = \mathbf{W}^* \mathbf{B}, \quad \hat{\mathbf{C}} = \mathbf{C} \mathbf{V},$$

where

$$\mathbf{W}^* = \begin{bmatrix} \mathbf{C}(s_{k+1} \mathbf{E} - \mathbf{A})^{-1} \\ \vdots \\ \mathbf{C}(s_{2k} \mathbf{E} - \mathbf{A})^{-1} \end{bmatrix}, \quad \mathbf{V} = [ (s_1 \mathbf{E} - \mathbf{A})^{-1} \mathbf{B} \quad \cdots \quad (s_k \mathbf{E} - \mathbf{A})^{-1} \mathbf{B} ]$$

(b) **Sylvester equations and projectors.** The solution of an appropriate Sylvester equation  $\mathbf{A}\mathbf{X} + \mathbf{X}\mathbf{H} + \mathbf{B}\mathbf{G} = \mathbf{0}$ , provides a projector that interpolates the original system  $\mathbf{C}, \mathbf{A}, \mathbf{B}$  at minus the eigenvalues of  $\mathbf{H}$ . Therefore the projectors above can be obtained by solving Sylvester equations. ■

# Outline

- 1 Krylov approximation methods
- 2 The Arnoldi and the Lanczos procedures
  - The Arnoldi procedure
  - The Lanczos procedure
  - An example
- 3 Krylov methods and moment matching
  - Remarks
- 4 Rational interpolation by Krylov projection
  - Realization by projection
  - Interpolation by projection
- 5 Choice of Krylov projection points: Optimal  $\mathcal{H}_2$  model reduction**
- 6 Summary: Lectures II and III

## Choice of Krylov projection points:

Optimal  $\mathcal{H}_2$  model reduction

Recall: the  $\mathcal{H}_2$  norm of a stable system is:

$$\|\Sigma\|_{\mathcal{H}_2} = \left( \int_{-\infty}^{+\infty} \mathbf{h}^2(t) dt \right)^{1/2}$$

where  $\mathbf{h}(t) = \mathbf{C}e^{\mathbf{A}t}\mathbf{B}$ ,  $t \geq 0$ , is the impulse response of  $\Sigma$ .

Goal: construct a *Krylov projector* such that

$$\Sigma_k = \arg \min_{\substack{\deg(\hat{\Sigma}) = r \\ \hat{\Sigma} : \text{stable}}} \|\Sigma - \hat{\Sigma}\|_{\mathcal{H}_2} = \left( \int_{-\infty}^{+\infty} (\mathbf{h} - \hat{\mathbf{h}})^2(t) dt \right)^{1/2}$$

## First-order necessary optimality conditions

Let  $(\hat{\mathbf{A}}, \hat{\mathbf{B}}, \hat{\mathbf{C}})$  solve the optimal  $\mathcal{H}_2$  problem and let  $\hat{\lambda}_i$  denote the eigenvalues of  $\hat{\mathbf{A}}$ . The **necessary conditions** are

$$\mathbf{H}(-\hat{\lambda}_i^*) = \hat{\mathbf{H}}(-\hat{\lambda}_i^*) \quad \text{and} \quad \left. \frac{d}{ds} \mathbf{H}(s) \right|_{s=-\hat{\lambda}_i^*} = \left. \frac{d}{ds} \hat{\mathbf{H}}(s) \right|_{s=-\hat{\lambda}_i^*}$$

Thus the reduced system has to match the first two moments of the original system at the *mirror images* of the eigenvalues of  $\hat{\mathbf{A}}$ .

**The  $\mathcal{H}_2$  norm:** if  $\mathbf{H}(s) = \sum_{k=1}^n \frac{\phi_k}{s-\lambda_k} \Rightarrow \boxed{\|\mathbf{H}\|_{\mathcal{H}_2}^2 = \sum_{k=1}^n c_k \mathbf{H}(-\lambda_k^*)}$

**Corollary.** With  $\hat{\mathbf{H}}(s) = \sum_{k=1}^r \frac{\hat{\phi}_k}{s-\hat{\lambda}_k}$ , the  $\mathcal{H}_2$  norm of the error system, is

$$\mathcal{J} = \|\mathbf{H} - \hat{\mathbf{H}}\|_{\mathcal{H}_2}^2 = \sum_{i=1}^n \phi_i [\mathbf{H}(-\lambda_i) - \hat{\mathbf{H}}(-\lambda_i)] + \sum_{j=1}^r \hat{\phi}_j [\hat{\mathbf{H}}(-\hat{\lambda}_j) - \mathbf{H}(-\hat{\lambda}_j)]$$

**Conclusion.** The  $\mathcal{H}_2$  error is due to the *mismatch* of the transfer functions  $\mathbf{H} - \hat{\mathbf{H}}$  at the *mirror images* of the full-order and reduced system poles  $\lambda_i, \hat{\lambda}_j$ .

## An iterative algorithm

Let the system obtained after the  $(j - 1)^{\text{st}}$  step be  $(\mathbf{C}_{j-1}, \mathbf{A}_{j-1}, \mathbf{B}_{j-1})$ , where  $\mathbf{A}_{j-1} \in \mathbb{R}^{k \times k}$ ,  $\mathbf{B}_{j-1}, \mathbf{C}_{j-1}^* \in \mathbb{R}^k$ . At the  $j^{\text{th}}$  step the system is obtained as follows

$$\mathbf{A}_j = (\mathbf{W}_j^* \mathbf{V}_j)^{-1} \mathbf{W}_j^* \mathbf{A} \mathbf{V}_j, \quad \mathbf{B}_j = (\mathbf{W}_j^* \mathbf{V}_j)^{-1} \mathbf{W}_j^* \mathbf{B}, \quad \mathbf{C}_j = \mathbf{C} \mathbf{V}_j,$$

where

$$\mathbf{V}_j = [(\lambda_1 \mathbf{I} - \mathbf{A})^{-1} \mathbf{B}, \dots, (\lambda_k \mathbf{I} - \mathbf{A})^{-1} \mathbf{B}],$$

$$\mathbf{W}_j^* = [\mathbf{C}(\lambda_1 \mathbf{I} - \mathbf{A})^{-1}, \dots, \mathbf{C}(\lambda_k \mathbf{I} - \mathbf{A})^{-1}],$$

and:  $-\lambda_1, \dots, -\lambda_k \in \sigma(\mathbf{A}_{j-1})$ ,

i.e.,  $-\lambda_i$  are the eigenvalues of the  $(j - 1)^{\text{st}}$  iterate  $\mathbf{A}_{j-1}$ .

The Newton step: can be computed explicitly

$$\begin{bmatrix} \lambda_1^{(k)} \\ \lambda_2^{(k)} \\ \vdots \end{bmatrix} \leftarrow \begin{bmatrix} \lambda_1^{(k)} \\ \lambda_2^{(k)} \\ \vdots \end{bmatrix} - \mathbf{J}^{-1} \begin{bmatrix} \lambda_1^{(k-1)} \\ \lambda_2^{(k-1)} \\ \vdots \end{bmatrix}$$

$\Rightarrow$  **local convergence guaranteed.**



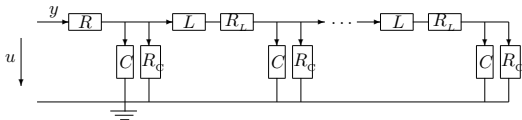
## An iterative rational Krylov algorithm (IRKA)

The proposed algorithm produces a reduced order model  $\hat{\mathbf{H}}(s)$  that satisfies the interpolation-based conditions, i.e.

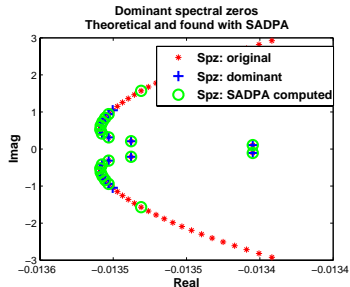
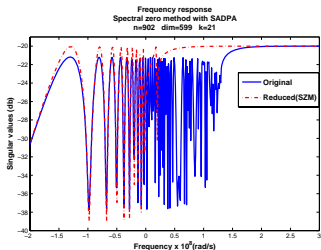
$$\mathbf{H}(-\hat{\lambda}_i^*) = \hat{\mathbf{H}}(-\hat{\lambda}_i^*) \quad \text{and} \quad \left. \frac{d}{ds} \mathbf{H}(s) \right|_{s=-\hat{\lambda}_i^*} = \left. \frac{d}{ds} \hat{\mathbf{H}}(s) \right|_{s=-\hat{\lambda}_i^*}$$

- 1 Make an initial selection of  $\sigma_i$ , for  $i = 1, \dots, k$
- 2  $\mathbf{W} = [(\sigma_1 \mathbf{I} - \mathbf{A}^*)^{-1} \mathbf{C}^*, \dots, (\sigma_k \mathbf{I} - \mathbf{A}^*)^{-1} \mathbf{C}^*]$
- 3  $\mathbf{V} = [(\sigma_1 \mathbf{I} - \mathbf{A})^{-1} \mathbf{B}, \dots, (\sigma_k \mathbf{I} - \mathbf{A})^{-1} \mathbf{B}]$
- 4 while (not converged)
  - $\hat{\mathbf{A}} = (\mathbf{W}^* \mathbf{V})^{-1} \mathbf{W}^* \mathbf{A} \mathbf{V}$ ,
  - $\sigma_i \leftarrow -\lambda_i(\hat{\mathbf{A}}) + \text{Newton correction}$ ,  $i = 1, \dots, k$
  - $\mathbf{W} = [(\sigma_1 \mathbf{I} - \mathbf{A}^*)^{-1} \mathbf{C}^*, \dots, (\sigma_k \mathbf{I} - \mathbf{A}^*)^{-1} \mathbf{C}^*]$
  - $\mathbf{V} = [(\sigma_1 \mathbf{I} - \mathbf{A})^{-1} \mathbf{B}, \dots, (\sigma_k \mathbf{I} - \mathbf{A})^{-1} \mathbf{B}]$
- 5  $\hat{\mathbf{A}} = (\mathbf{W}^* \mathbf{V})^{-1} \mathbf{W}^* \mathbf{A} \mathbf{V}$ ,  $\hat{\mathbf{B}} = (\mathbf{W}^* \mathbf{V})^{-1} \mathbf{W}^* \mathbf{B}$ ,  $\hat{\mathbf{C}} = \mathbf{C} \mathbf{V}$

## Moderate-dimensional example



- total system variables  $n = 902$ , independent variables  $\dim = 599$ , reduced dimension  $k = 21$
- reduced model captures dominant modes



$\mathcal{H}_\infty$  and  $\mathcal{H}_2$  error norms

## Relative norms of the error systems

Reduction Method $n = 902, \dim = 599, k = 21$	$\mathcal{H}_\infty$	$\mathcal{H}_2$
PRIMA	1.4775	-
Spectral Zero Method with SADPA	0.9628	0.841
Optimal $\mathcal{H}_2$	<b>0.5943</b>	<b>0.4621</b>
Balanced truncation (BT)	0.9393	0.6466
Riccati Balanced Truncation (PRBT)	0.9617	0.8164

# Outline

- 1 Krylov approximation methods
- 2 The Arnoldi and the Lanczos procedures
  - The Arnoldi procedure
  - The Lanczos procedure
  - An example
- 3 Krylov methods and moment matching
  - Remarks
- 4 Rational interpolation by Krylov projection
  - Realization by projection
  - Interpolation by projection
- 5 Choice of Krylov projection points: Optimal  $\mathcal{H}_2$  model reduction
- 6 **Summary: Lectures II and III**

## Approximation methods: Summary

## Krylov

- Realization
- Interpolation
- Lanczos
- Arnoldi

## SVD

## Nonlinear systems

- POD methods
- Empirical Gramians

## Linear systems

- Balanced truncation
- Hankel approximation

## Properties

- numerical efficiency
- $n \gg 10^3$
- choice of matching moments

## Krylov/SVD Methods

## Properties

- Stability
- Error bound
- $n \approx 10^3$

## Complexity considerations

### • Dense problems

#### Major cost **Balanced Truncation**:

Compute gramians  $\approx 30N^3$  (eigenvalue decomp.)

Perform balancing  $\approx 25N^3$  (sing. value decomp.)

#### **Rational Krylov approximation**:

Decompose  $(\mathbf{A} - \sigma_j \mathbf{E})$  for  $k$  points  $\approx \frac{2}{3} kN^3$

**Remark** : Iterations (Sign, Smith) can accelerate the computation of gramians (esp. on parallel machines)

### • Approximate and/or sparse decompositions

#### Major cost **Balanced Truncation**:

Compute gramians  $\approx c_1 \alpha kN$

Perform balancing  $O(n^3)$

#### **Rational Krylov approximation**:

Iterative solves for  $(\mathbf{A} - \sigma_j \mathbf{E})\mathbf{x} = \mathbf{b} \approx c_2 k \alpha N$ , where  $k$  = number of expansion points;  $\alpha$  = average number of non-zero elements per row in  $\mathbf{A}$ ,  $\mathbf{E}$ .